## Efficient compilation of complex tensor algebra expressions

Martin Sandve Alnæs
Center for Biomedical Computing
[ simula . research laboratory ]

- by thinking constantly about it

June 5th, FEniCS 2012

# UFL is a DSL, symbolic backend, and frontend to form compilers 

- Users recently reported scalabiliy problems when compiling complicated equations
- After some profiling sessions I reduced the memory usage by a factor 10 for one case
- Next I have made an attempt at faster form compiler algorithms, which I will show

UFL expressions are represented in a symbolic expression tree

IntValue

## CellVolume

## Coefficient

$$
\begin{gathered}
\mathrm{e}=\text { dot(as_vector }((1, \text { cell.volume })), \\
\text { Coefficient(V) })
\end{gathered}
$$

## Some quick design points

- Expr objects are immutable for easy sharing
- Conservative approach to automatic simplifications
- Canonical ordering of sum and product:



# UFL simplifies some expressions on construction 


$0+f$
op(c1)


## Simplifications critical during differentiation algorithm

as_tensor(A[i,j], (i,j))


- $d / d x\left(x^{*} g(y)\right)=1^{*} g+x^{*} 0->g$


## Performance must scale as $O(n)$ with size of expression

- This means almost anything must be $\mathrm{O}(1)$
- In particular __eq__ and __hash__!


## Transformations must be safe for floating point computations

- Def eps: 1 + eps > 1
- $(1+e p s / 2)+e p s / 2==1$
- $1+(e p s / 2+e p s / 2)>1$


## I will take this expression through the compiler algorithms

a, b, c = scalar coefficients

$$
\mathrm{u}=\operatorname{as} \_\operatorname{vector}((0, a, b))
$$

$$
\mathrm{v}=\mathrm{as} \_\operatorname{vector}((\mathrm{c}, \mathrm{~b}, \mathrm{a}))
$$

$$
e=\operatorname{dot}(u, v)
$$

Anticipate result:

$$
\begin{aligned}
t & =a^{*} b \\
e & =t+t
\end{aligned}
$$

The expression tree after translating dot to index notation


## Placing nodes in array

| Index | $\mathrm{V}[\mathrm{i}]$ | Shape | Size |
| :---: | :---: | :---: | :---: |
| 0 | 0 | ()$+()$ | 1 |
| 1 | a | ()$+()$ | 1 |
| 2 | b | ()$+()$ | 1 |
| 3 | c | ()$+()$ | 1 |
| 4 | $<0, \mathrm{a}, \mathrm{b}>$ | $(3)+()$, | 3 |
| 5 | $<\mathrm{c}, \mathrm{b}, \mathrm{a}>$ | $(3)+()$, | 3 |
| 6 | $\mathrm{u}[i]$ | ()$+(3)$, | 3 |
| 7 | $\mathrm{v}[i]$ | ()$+(3)$, | 3 |
| 8 | $\mathrm{u}[\mathrm{i}]^{\star} \mathrm{v}[i]$ | ()$+(3)$, | 3 |
| 9 | ISum(V$[8], \mathrm{i})$ | ()$+()$ | 1 |

Scalar subexpressions are assigned unique value numbers

| Index | $\mathrm{V}[\mathrm{i}]$ | Value number |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 1 | a | 1 |
| 2 | b | 2 |
| 3 | c | 3 |
| 4 | $<0, \mathrm{a}, \mathrm{b}>$ | $0,1,2$ |
| 5 | $<\mathrm{c}, \mathrm{b}, \mathrm{a}>$ | $3,2,1$ |
| 6 | $\mathrm{u}[\mathrm{i}]$ | $0,1,2$ |
| 7 | $\mathrm{v}[\mathrm{i}]$ | $3,2,1$ |
| 8 | $\mathrm{u}[\mathrm{i}]^{*} \mathrm{~V}[\mathrm{i}]$ | $4,5,6$ |
| 9 | Sum $(V[8], \mathrm{i})$ | 7 |

## Scalar subexpressions are reevaluated and placed in a new array

| Index | $S[1]$ | Simplifies to |
| :---: | :---: | :---: |
| 0 | 0 |  |
| 1 | $a$ |  |
| 2 | b |  |
| 3 | c |  |
| 4 | $\mathrm{~S}[0]^{*} \mathrm{~S}[3]$ | $0^{*} \mathrm{c}=0$ |
| 5 | S[]$^{*} \mathrm{~S}[3]$ | $\mathrm{a}^{*} \mathrm{~b}$ |
| 6 | $\mathrm{~S}[0]^{*} \mathrm{~S}[3]$ | $\mathrm{b}^{*} \mathrm{a}=\mathrm{a}^{*} \mathrm{~b}$ |
| 7 | $\mathrm{~S}[4]+\mathrm{S}[5]+\mathrm{S}[6]$ | $\mathrm{a}^{*} \mathrm{~b}+\mathrm{a}^{*} \mathrm{~b}$ |

## Throwing away the array only keeping the final expression

## Placing nodes in array!

| Index | $V[i]$ | Shape | Size |
| :---: | :---: | :---: | :---: |
| 0 | $a$ | ()$+()$ | 1 |
| 1 | $b$ | ()$+()$ | 1 |
| 2 | $a^{*} b$ | ()$+()$ | 1 |
| 3 | $a^{*} b+a^{*} b$ | ()$+()$ | 1 |

## Analyzing dependencies

Index
V[i]
Dep.
Rev. Dep.
()
()
$(0,1)$
$(3,3)$

3
$a * b+a * b$
$(2,2)$
()
(2,)
(2,)

## Final steps in compiler

- Partition final array by dependencies on $\mathrm{x}, \mathrm{u}, \mathrm{v}$
- Heuristically pick best candidates for subexpressions to place in intermediate variables in generated code
- Format expressions and assignment statements within nested loops
- FEM library specific code generation in separate plugin class, e.g. how to evaluate geometry and coefficients, how to loop over quadrature points and basis functions


## Outlook

- bzr branch Ip:uflacs
- Can generate dolfin::Expression classes
- Soon SFC can use uflacs to compile forms
- Want to merge algorithms into FFC
- Write plugin class to compile to other FEM libs

